

# Lecture 9: Procedures

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601.229 Computer Systems Fundamentals



# Control flow (part 2)

- ▶ Procedures
- ▶ Stacks:
  - ▶ Procedure calls and returns
  - ▶ Storage for local variables and temporary values
- ▶ Today's example programs are linked as `control2.zip` on the course website

# Procedures

# Procedures, call stack

- ▶ Procedures (a.k.a. functions, subroutines), the most important abstraction in programming
  - ▶ Can you imagine trying to write programs without them?
- ▶ *Call stack*: hardware-supported, runtime data structure
  - ▶ Stores *return addresses* so procedures know where to return to
  - ▶ Used to allocate *stack frames*: per-procedure-call storage area for local variables, temporary values, and (sometimes) argument values
  - ▶ As name suggests, is a stack, LIFO discipline (push and pop)

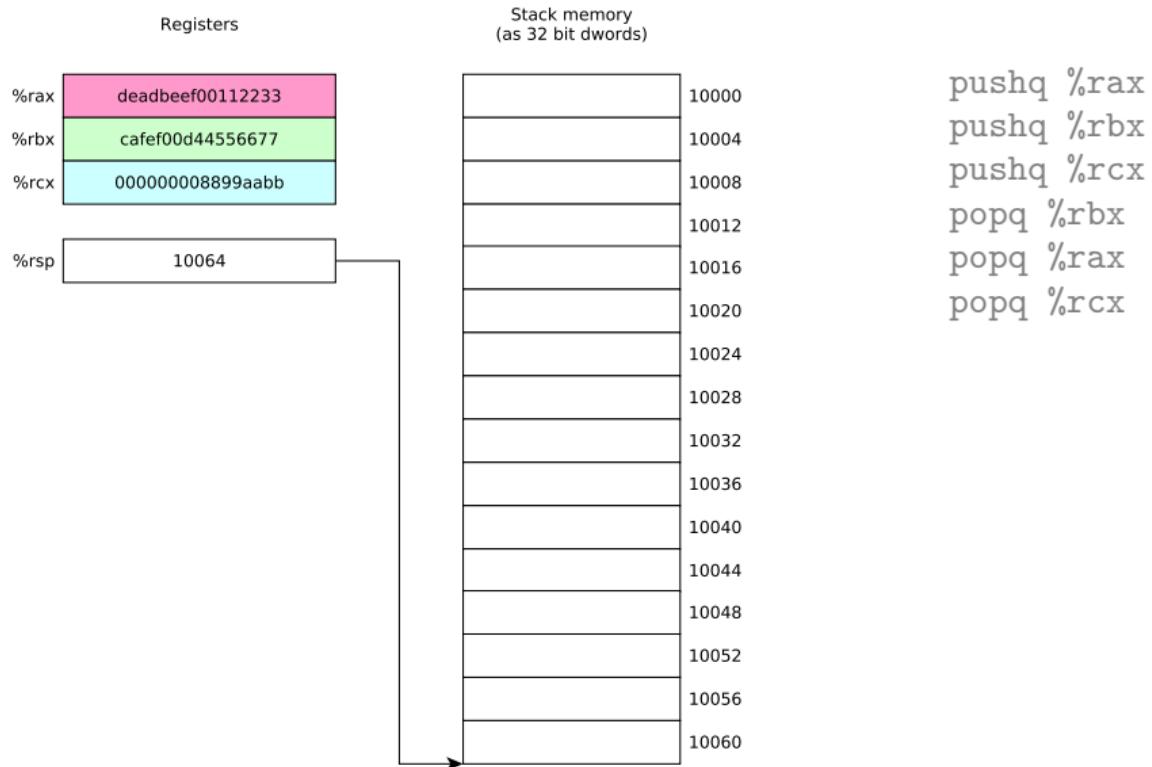
# Stack pointer, instruction pointer

- ▶ *Stack pointer* register `%rsp`: contains address of current “top” of stack
  - ▶ Important: stack grows towards lower addresses, so top of stack is at lower address than bottom of stack
- ▶ *Instruction pointer* register `%rip`: contains code address of next instruction to be updated
  - ▶ Control flow changes the value of `%rip`
- ▶ Other architectures use the name “program counter” rather than “instruction pointer”, but they’re the same thing

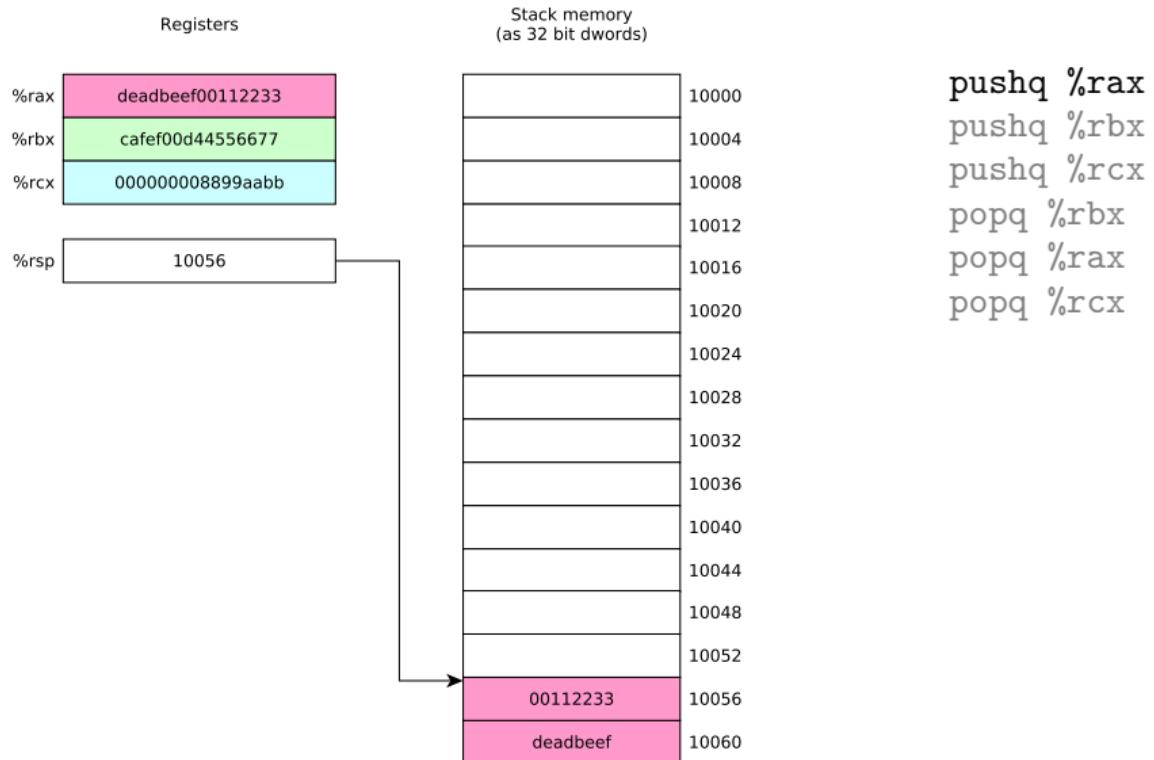
# push and pop

- ▶ push: push a data value onto the call stack
  - ▶ E.g., `pushq %rax`
    - ▶ Decrement `%rsp` by 8
    - ▶ Store value in `%rax` at memory location pointed-to by `%rsp`
- ▶ pop: pop a data value from the call stack
  - ▶ E.g., `popq %rax`
    - ▶ Load value at memory location pointed-to by `%rsp` into `%rax`
    - ▶ Increment `%rsp` by 8
- ▶ push and pop are amazingly useful for saving and restoring register values
- ▶ Various size operands (1, 2, 4, 8 bytes) can be pushed and popped; need to consider alignment

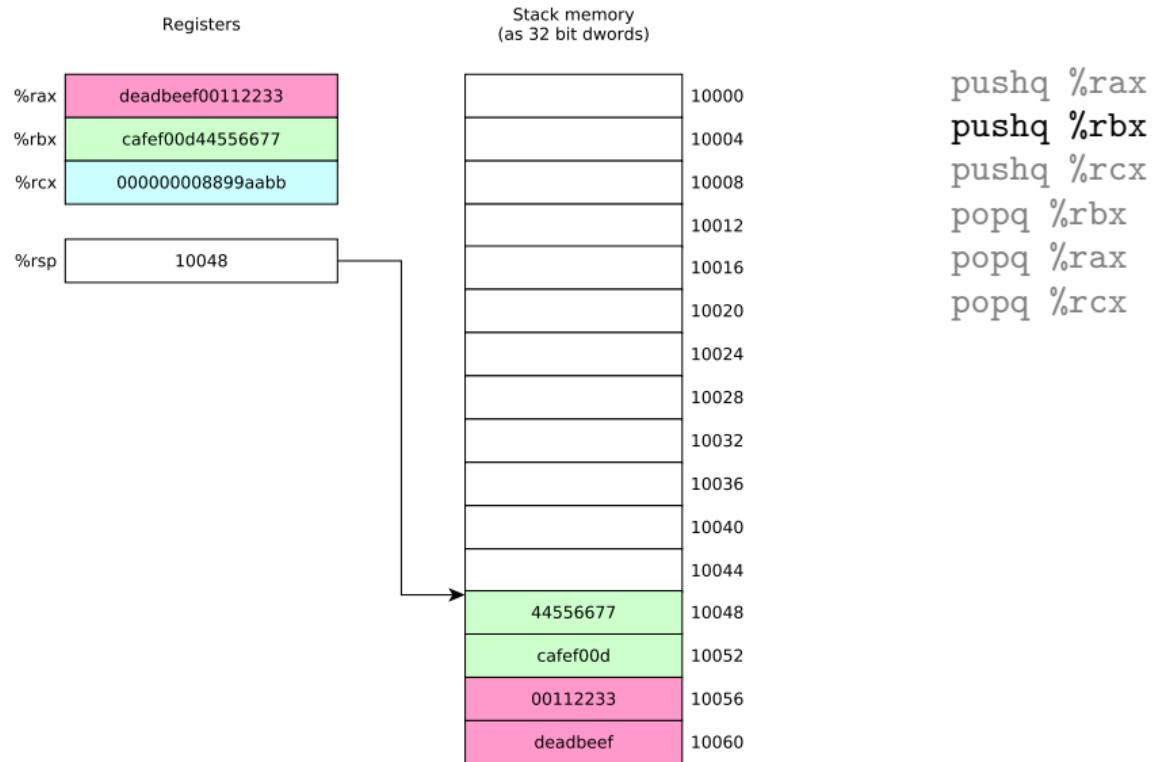
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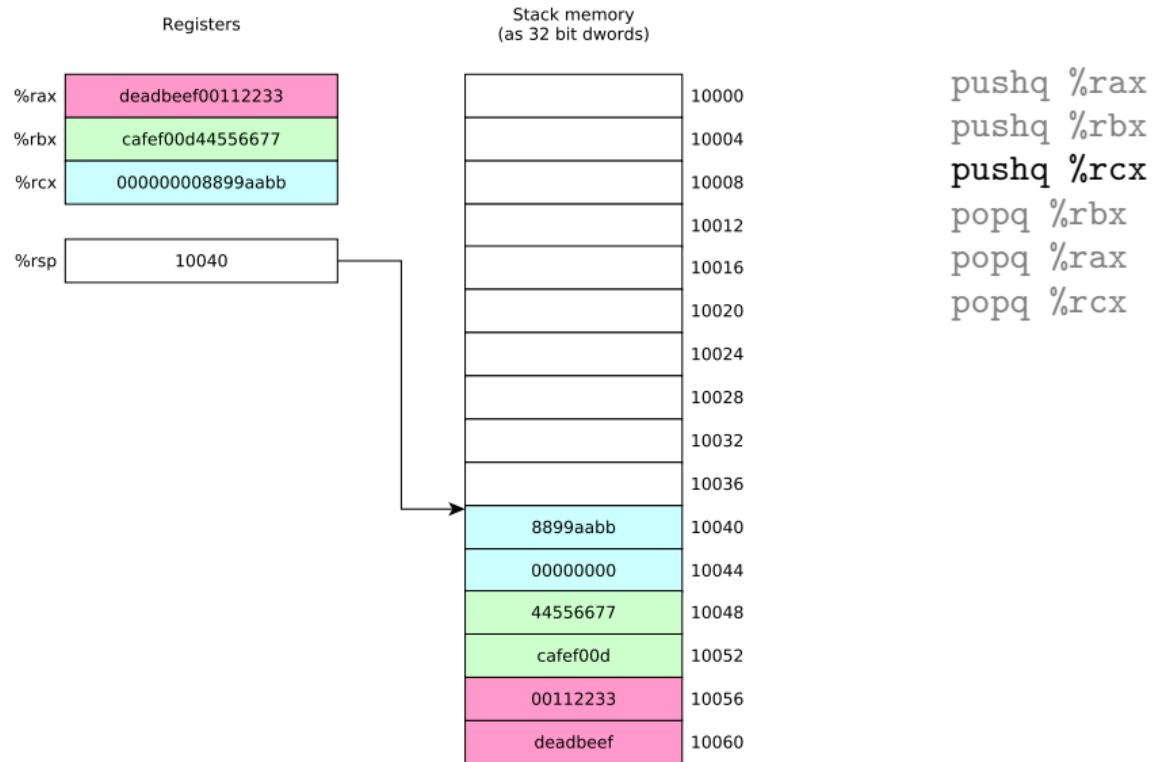
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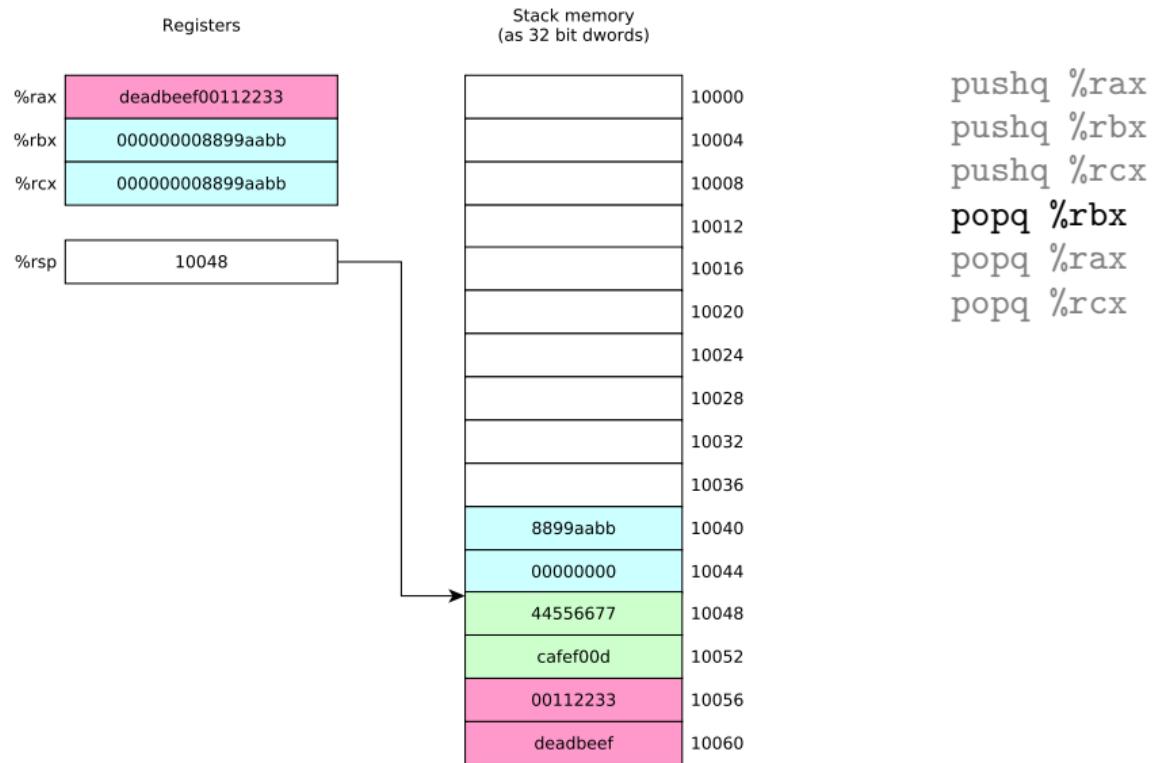
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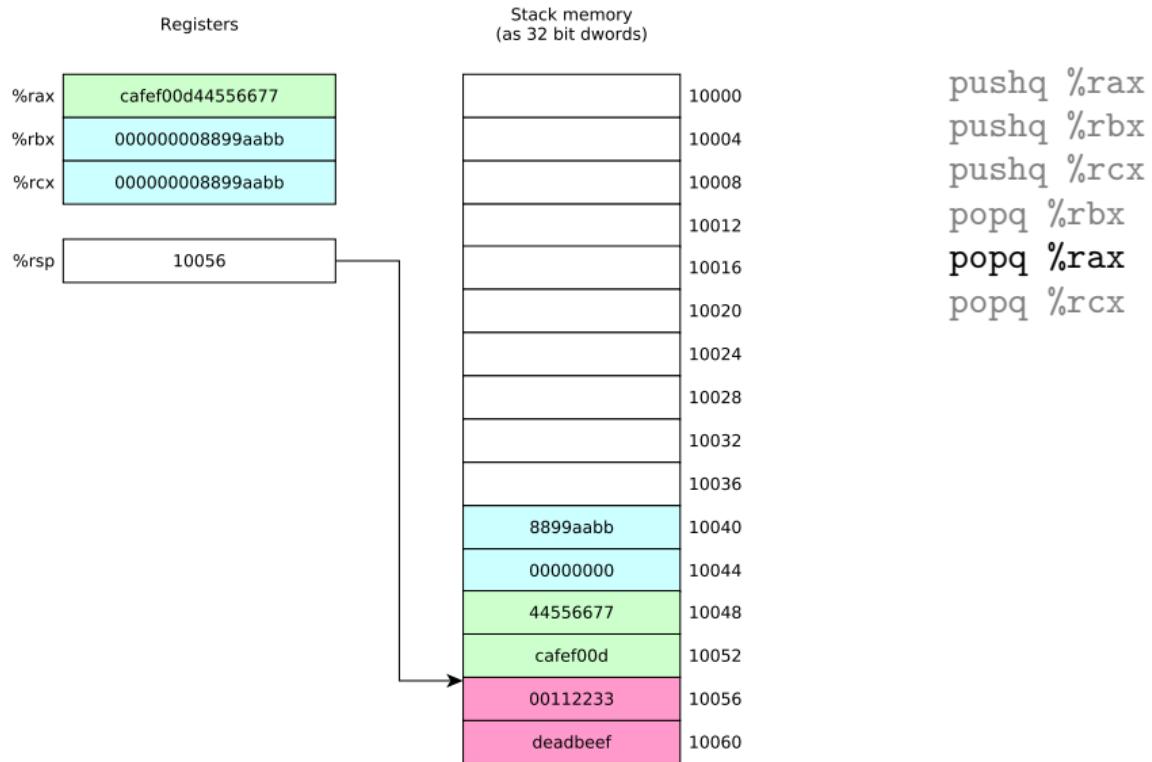
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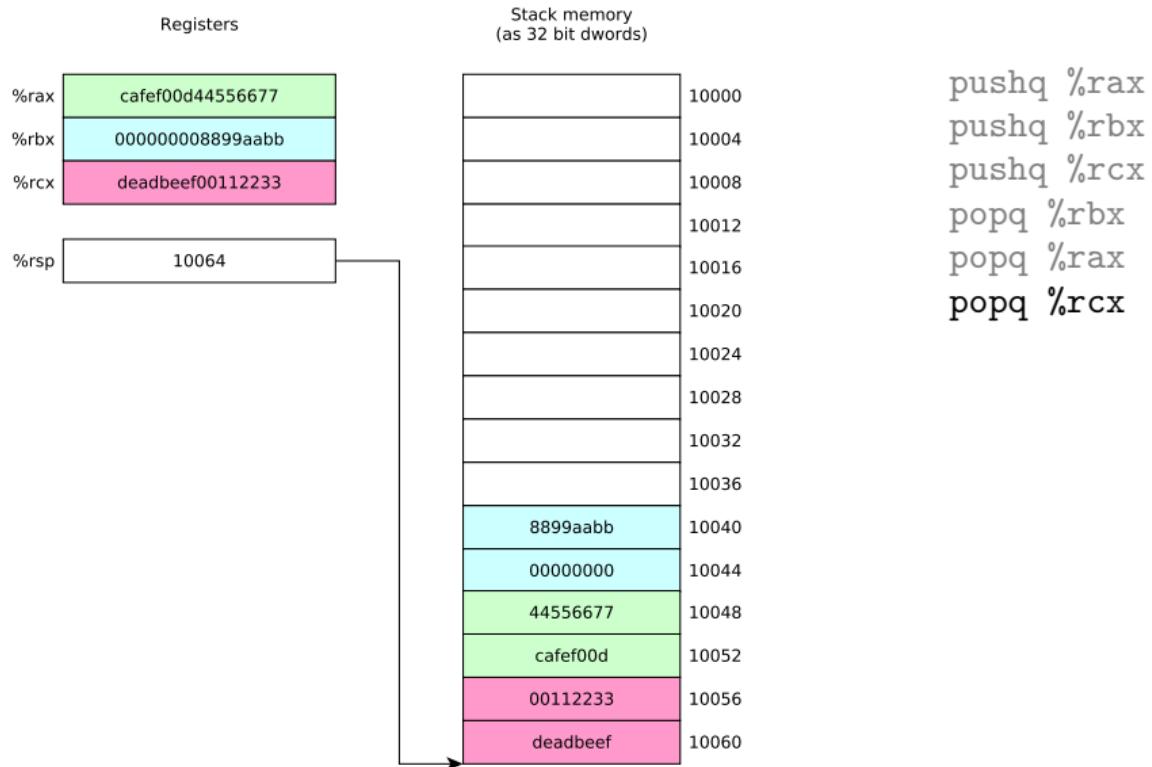
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# call and ret

- ▶ call instruction: calls procedure
  - ▶ %rip contains address of instruction following call instruction
  - ▶ Push %rip onto stack (as though pushq %rip was executed): this is the *return address*
  - ▶ Change %rip to address of first instruction of called procedure
  - ▶ Called procedure starts executing
- ▶ ret instruction: return from procedure
  - ▶ Pop saved return address from stack into %rip (as though popq %rip was executed)
  - ▶ Execution continues at return address

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- ▶ E.g., storage for an 8 byte value should be stored at an address which is a multiple of 8
- ▶ This is true of stack-allocated values!
- ▶ The Linux x86-64 calling conventions require `%rsp` to be a multiple of 16 at the point of a procedure call (to ensure that 16 byte values can be accessed on the stack if necessary)
- ▶ **Issue:** on entry to a procedure,  $\%rsp \bmod 16 = 8$  because the `call` instruction (which called the procedure) pushed `%rip` (the program counter) onto the stack

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- ▶ The Linux `printf` function will segfault if the stack is misaligned

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  - ▶ They allow your code to interoperate with other code, including library routines and (OS) system calls
- ▶ **Always follow the appropriate register use conventions**

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- ▶ Callee-saved registers: %rbx, %rbp, %r12, %r13, %14, %r15

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- ▶ *Caller-saved* registers: caller must *not* assume that the procedure call will preserve their value
  - ▶ In general any procedure can freely modify them
  - ▶ A caller might need to save their contents to memory prior to calling a procedure and restore the value afterwards

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  - ▶ Use callee-saved registers for longer term values that need to persist across procedure calls
  - ▶ Use pushq/popq to save and restore their values on procedure entry and exit

# Recursive Fibonacci computation

Compute  $n$ th Fibonacci number recursively (warning: exponential-time algorithm!)

The call stack inherently allows recursion: there is nothing special we need to do to make it work

Recall that

$$fib(0) = 0$$

$$fib(1) = 1$$

$$\text{For } n > 1, fib(n) = fib(n - 2) + fib(n - 1)$$

# Recursive Fibonacci function (see fibRec.S for full program)

```
fib:  
    cmpl $2, %edi  
    jae .LrecursiveCase  
    movl %edi, %eax  
    ret  
.LrecursiveCase:  
    /* recursive case */  
    pushq %r12  
    movl %edi, %r12d  
    subl $2, %edi  
    call fib  
    movl %r12d, %edi  
    subl $1, %edi  
    movl %eax, %r12d  
    call fib  
    addl %r12d, %eax  
    popq %r12  
    ret  
    /* check base case */  
    /* if n>=2, do recursive case */  
    /* base case, just return n */  
    /* preserve value of %r12 */  
    /* save n in %r12 */  
    /* compute n-2 */  
    /* compute fib(n-2) */  
    /* put saved n in %edi */  
    /* compute n-1 */  
    /* save fib(n-2) in %r12 */  
    /* compute fib(n-1) */  
    /* return fib(n-2)+fib(n-1) */  
    /* restore value of %r12 */  
    /* done */
```

# Running the program (with N=9)

```
$ gcc -c -g -no-pie -o fibRec.o fibRec.S
$ gcc -no-pie -o fibRec fibRec.o
$ ./fibRec
fib(9) = 34
```

# Clicker quiz!

Clicker quiz omitted from public slides

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- ▶ Could use heap allocation (i.e., malloc, free)
  - ▶ Has overhead due to bookkeeping, locking
- ▶ The call stack is an ideal place to allocate storage for local variables

# Stack allocation

- ▶ Stack allocation of storage is simple:
  - ▶ To allocate  $n$  bytes, subtract  $n$  from `%rsp`
    - ▶ Updated `%rsp` is a pointer to the beginning of the allocated memory
  - ▶ To deallocate  $n$  bytes, add  $n$  to `%rsp`
- ▶ Complication: instructions such as `push` and `pop` change `%rsp`
- ▶ Solution: use the *frame pointer* register `%rbp` to keep track of allocated memory area

# Using the frame pointer

On entry to procedure:

```
pushq %rbp  
movq %rsp, %rbp  
subq $N, %rsp
```

Before returning from procedure:

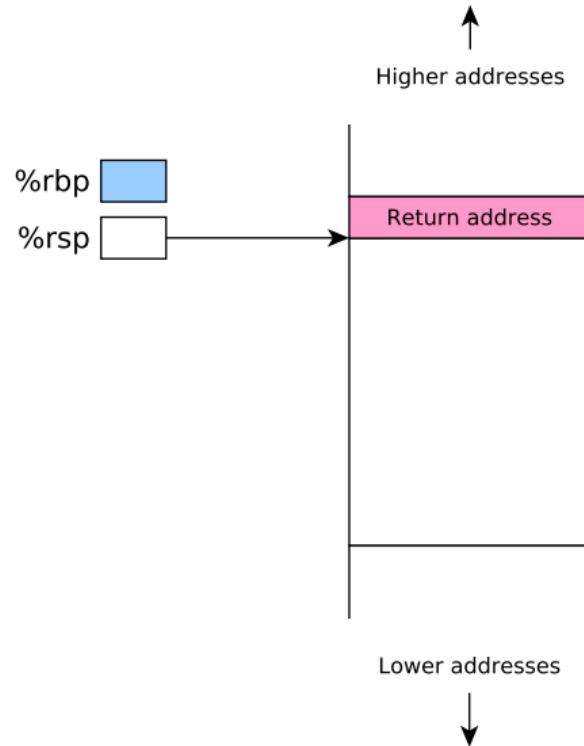
```
addq $N, %rsp  
popq %rbp
```

`%rbp` points to a memory location *just above* a block of  $N$  bytes allocated in the current stack frame. Note that

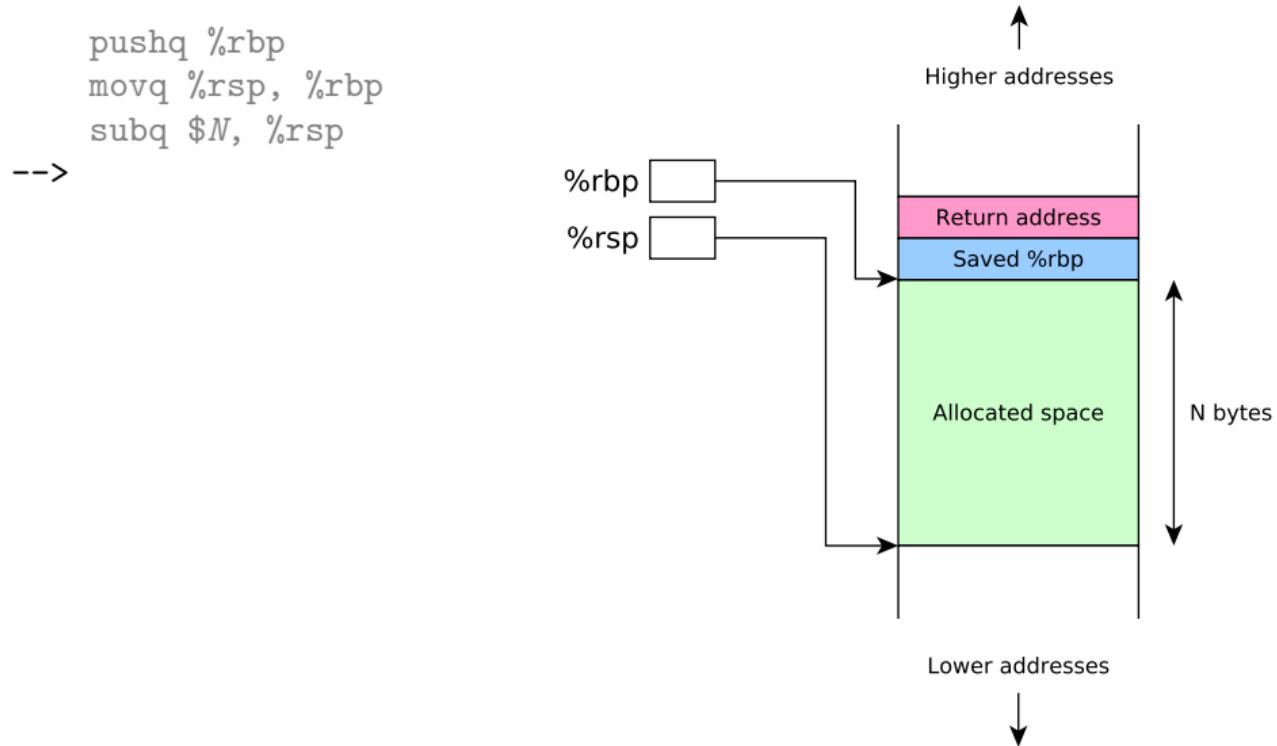
- ▶  $N$  should be a multiple of 16 to ensure correct stack alignment
- ▶ The function will access memory locations in the allocated block using *negative* offsets from `%rbp`

# Before allocating space in stack frame

```
--> pushq %rbp  
    movq %rsp, %rbp  
    subq $N, %rsp
```



# After allocating space in stack frame



# Putting it all together

- ▶ Let's examine a simple program which
  - ▶ Reads two 64 bit integer values from user
  - ▶ Computes their sum using a function
  - ▶ Prints out the sum
- ▶ Calling `scanf` to read input requires variables in which to store input values: we'll allocate them on the stack

# addLongs, C version

```
#include <stdio.h>

long addLongs(long a, long b);

int main(void) {
    long x, y, sum;
    printf("Enter two integers: ");
    scanf("%ld %ld", &x, &y);
    sum = addLongs(x, y);
    printf("Sum is %ld\n", sum);
}

long addLongs(long a, long b) {
    return a + b;
}
```

# addLongs, assembly version

```
.section .rodata
sPromptMsg: .string "Enter two integers: "
sInputFmt: .string "%ld %ld"
sResultMsg: .string "Sum is %ld\n"

.section .text
.globl main
.align 16
main:
    pushq %rbp
    movq %rsp, %rbp
    subq $16, %rsp
    movl $0, %eax
    movq $sPromptMsg, %rdi
    call printf

    movl $0, %eax
    movq $sInputFmt, %rdi
    leaq -16(%rbp), %rsi
    leaq -8(%rbp), %rdx
    call scanf

    movq -16(%rbp), %rdi
    movq -8(%rbp), %rsi
    call addLongs

    movq $sResultMsg, %rdi
    movq %rax, %rsi
    call printf

    addq $16, %rsp
    popq %rbp
    ret

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main:
    pushq %rbp      !-- save orig value of %rbp
    movq %rsp, %rbp
    subq $16, %rsp

    movl $0, %eax
    movq $sPromptMsg, %rdi
    call printf

    movl $0, %eax
    movq $sInputFmt, %rdi
    leaq -16(%rbp), %rsi
    leaq -8(%rbp), %rdx
    call scanf

    movq -16(%rbp), %rdi
    movq -8(%rbp), %rsi
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    movq $sResultMsg, %rdi
    movq %rax, %rsi
    call printf

    addq $16, %rsp
    popq %rbp
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    pushq %rbp
    movq %rsp, %rbp    <-- %rbp points to top
    subq $16, %rsp     of alloc'ed area
    movl $0, %eax
    movq $sPromptMsg, %rdi
    call printf

    movl $0, %eax
    movq $sInputFmt, %rdi
    leaq -16(%rbp), %rsi
    leaq -8(%rbp), %rdx
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    movq $sResultMsg, %rdi
    movq %rax, %rsi
    call printf

    addq $16, %rsp
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    movq %rsp, %rbp
    subq $16, %rsp
    movl $0, %eax
    movq $sPromptMsg, %rdi
    call printf
    movl $0, %eax
    movq $sInputFmt, %rdi
    leaq -16(%rbp), %rsi  <-- pass address of 1st var
    leaq -8(%rbp), %rdx
    call scanf
    movq -16(%rbp), %rdi
    movq -8(%rbp), %rsi
    call addLongs
    movq $sResultMsg, %rdi
    movq %rax, %rsi
    call printf
    addq $16, %rsp
    popq %rbp
    ret
.align 16
addLongs:
    movq %rdi, %rax
    addq %rsi, %rax
    ret
```

# addLongs, assembly version

```
.section .rodata
sPromptMsg: .string "Enter two integers: "
sInputFmt: .string "%ld %ld"
sResultMsg: .string "Sum is %ld\n"

.section .text
.globl main
.align 16
main:
    pushq %rbp
    movq %rsp, %rbp
    subq $16, %rsp
    movl $0, %eax
    movq $sPromptMsg, %rdi
    call printf
    movl $0, %eax
    movq $sInputFmt, %rdi
    leaq -16(%rbp), %rsi
    leaq -8(%rbp), %rdx    <-- pass address of 2nd var
    call scanf
    movq -16(%rbp), %rdi
    movq -8(%rbp), %rsi
    call addLongs
    movq $sResultMsg, %rdi
    movq %rax, %rsi
    call printf
    addq $16, %rsp
    popq %rbp
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    movq %rdi, %rax
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    movq $sPromptMsg, %rdi
    call printf

    movl $0, %eax
    movq $sInputFmt, %rdi
    leaq -16(%rbp), %rsi
    leaq -8(%rbp), %rdx
    call scanf

    movq -16(%rbp), %rdi  <-- pass value of 1st var
    movq -8(%rbp), %rsi
    call addLongs

    movq $sResultMsg, %rdi
    movq %rax, %rsi
    call printf

    addq $16, %rsp
    popq %rbp
    ret

.align 16
addLongs:
    movq %rdi, %rax
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    subq $16, %rsp

    movl $0, %eax
    movq $sPromptMsg, %rdi
    call printf

    movl $0, %eax
    movq $sInputFmt, %rdi
    leaq -16(%rbp), %rsi
    leaq -8(%rbp), %rdx
    call scanf

    movq -16(%rbp), %rdi
    movq -8(%rbp), %rsi    <-- pass value of 2nd var
    call addLongs

    movq $sResultMsg, %rdi
    movq %rax, %rsi
    call printf

    addq $16, %rsp
    popq %rbp
    ret

.align 16
addLongs:
    movq %rdi, %rax
    addq %rsi, %rax
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```

# addLongs, assembly version

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.section .text
.globl main
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main:
    pushq %rbp
    movq %rsp, %rbp
    subq $16, %rsp

    movl $0, %eax
    movq $sPromptMsg, %rdi
    call printf

    movl $0, %eax
    movq $sInputFmt, %rdi
    leaq -16(%rbp), %rsi
    leaq -8(%rbp), %rdx
    call scanf

    movq -16(%rbp), %rdi
    movq -8(%rbp), %rsi
    call addLongs

    movq $sResultMsg, %rdi
    movq %rax, %rsi
    call printf

    addq $16, %rsp  <-- deallocate alloc'ed area
    popq %rbp
    ret

.align 16
addLongs:
    movq %rdi, %rax
    addq %rsi, %rax
    ret
```

# addLongs, assembly version

```
.section .rodata
sPromptMsg: .string "Enter two integers: "
sInputFmt: .string "%ld %ld"
sResultMsg: .string "Sum is %ld\n"

.section .text
.globl main
.align 16
main:
    pushq %rbp
    movq %rsp, %rbp
    subq $16, %rsp

    movl $0, %eax
    movq $sPromptMsg, %rdi
    call printf

    movl $0, %eax
    movq $sInputFmt, %rdi
    leaq -16(%rbp), %rsi
    leaq -8(%rbp), %rdx
    call scanf

    movq -16(%rbp), %rdi
    movq -8(%rbp), %rsi
    call addLongs

    movq $sResultMsg, %rdi
    movq %rax, %rsi
    call printf

    addq $16, %rsp
    popq %rbp          <-- restore orig value of %rbp
    ret

.align 16
addLongs:
    movq %rdi, %rax
    addq %rsi, %rax
    ret
```

# addLongs, assembly version

```
.section .rodata
sPromptMsg: .string "Enter two integers: "
sInputFmt: .string "%ld %ld"
sResultMsg: .string "Sum is %ld\n"

.section .text
.globl main
.align 16
main:
    pushq %rbp
    movq %rsp, %rbp
    subq $16, %rsp

    movl $0, %eax
    movq $sPromptMsg, %rdi
    call printf

    movl $0, %eax
    movq $sInputFmt, %rdi
    leaq -16(%rbp), %rsi
    leaq -8(%rbp), %rdx
    call scanf

    movq -16(%rbp), %rdi
    movq -8(%rbp), %rsi
    call addLongs

    movq $sResultMsg, %rdi
    movq %rax, %rsi
    call printf

    addq $16, %rsp
    popq %rbp
    ret

.align 16
addLongs: <-- does not use stack, ignore alignment :-P
    movq %rdi, %rax
    addq %rsi, %rax
    ret
```

# Running the program

```
$ gcc -c -no-pie -o addLongs.o addLongs.S
$ gcc -no-pie -o addLongs addLongs.o
$ ./addLongs
Enter two integers: 2 3
Sum is 5
```

# Running the program in gdb

```
$ gdb addLongs
...output omitted...
(gdb) break addLongs.S:28
Breakpoint 1 at 0x401172: file addLongs.S, line 28.
(gdb) run
Starting program: /home/daveho/.../src/control2/addLongs
Enter two integers: 3 4

Breakpoint 1, main () at addLongs.S:28
28          movq -16(%rbp), %rdi      /* pass first value */
(gdb) print *(long *)($rbp-16)
$1 = 3
(gdb) print *(long *)($rbp-8)
$2 = 4
```

# Running the program in gdb

```
$ gdb addLongs
...output omitted...
(gdb) break addLongs.S:28  <-- set breakpoint just after scanf returns
Breakpoint 1 at 0x401172: file addLongs.S, line 28.
(gdb) run
Starting program: /home/daveho/.../src/control2/addLongs
Enter two integers: 3 4

Breakpoint 1, main () at addLongs.S:28
28          movq -16(%rbp), %rdi      /* pass first value */
(gdb) print *(long *)($rbp-16)
$1 = 3
(gdb) print *(long *)($rbp-8)
$2 = 4
```

# Running the program in gdb

```
$ gdb addLongs
...output omitted...
(gdb) break addLongs.S:28
Breakpoint 1 at 0x401172: file addLongs.S, line 28.
(gdb) run                  <-- start the program
Starting program: /home/daveho/.../src/control2/addLongs
Enter two integers: 3 4

Breakpoint 1, main () at addLongs.S:28
28          movq -16(%rbp), %rdi      /* pass first value */
(gdb) print *(long *)($rbp-16)
$1 = 3
(gdb) print *(long *)($rbp-8)
$2 = 4
```

# Running the program in gdb

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(gdb) break addLongs.S:28
Breakpoint 1 at 0x401172: file addLongs.S, line 28.
(gdb) run
Starting program: /home/daveho/.../src/control2/addLongs
Enter two integers: 3 4    <-- enter input values

Breakpoint 1, main () at addLongs.S:28
28          movq -16(%rbp), %rdi      /* pass first value */
(gdb) print *(long *)($rbp-16)
$1 = 3
(gdb) print *(long *)($rbp-8)
$2 = 4
```

# Running the program in gdb

```
$ gdb addLongs
...output omitted...
(gdb) break addLongs.S:28
Breakpoint 1 at 0x401172: file addLongs.S, line 28.
(gdb) run
Starting program: /home/daveho/.../src/control2/addLongs
Enter two integers: 3 4

Breakpoint 1, main () at addLongs.S:28
28          movq -16(%rbp), %rdi      /* pass first value */
(gdb) print *(long *)($rbp-16)  <-- print first input value at -16(%rbp)
$1 = 3
(gdb) print *(long *)($rbp-8)
$2 = 4
```

# Running the program in gdb

```
$ gdb addLongs
...output omitted...
(gdb) break addLongs.S:28
Breakpoint 1 at 0x401172: file addLongs.S, line 28.
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Starting program: /home/daveho/.../src/control2/addLongs
Enter two integers: 3 4

Breakpoint 1, main () at addLongs.S:28
28          movq -16(%rbp), %rdi      /* pass first value */
(gdb) print *(long *)($rbp-16)
$1 = 3
(gdb) print *(long *)($rbp-8)    <-- print second input value at -8(%rbp)
$2 = 4
```

# Running the program in gdb

```
$ gdb addLongs
...output omitted...
(gdb) break addLongs.S:28
Breakpoint 1 at 0x401172: file addLongs.S, line 28.
(gdb) run
Starting program: /home/daveho/.../src/control2/addLongs
Enter two integers: 3 4

Breakpoint 1, main () at addLongs.S:28
28          movq -16(%rbp), %rdi      /* pass first value */
(gdb) print *(long *)($rbp-16)
$1 = 3
(gdb) print *(long *)($rbp-8)
$2 = 4
```