



# Lecture 9: Procedures

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# Control flow (part 2)

- Procedures
- Stacks:
  - Procedure calls and returns
  - Storage for local variables and temporary values
- Today's example programs are linked as `control2.zip` on the course website



# Procedures

# Procedures, call stack

- Procedures (a.k.a. functions, subroutines), the most important abstraction in programming
  - Can you imagine trying to write programs without them?
- *Call stack*: hardware-supported, runtime data structure
  - Stores *return addresses* so procedures know where to return to
  - Used to allocate *stack frames*: per-procedure-call storage area for local variables, temporary values, and (sometimes) argument values
  - As name suggests, is a stack, LIFO discipline (push and pop)



# Stack pointer, instruction pointer

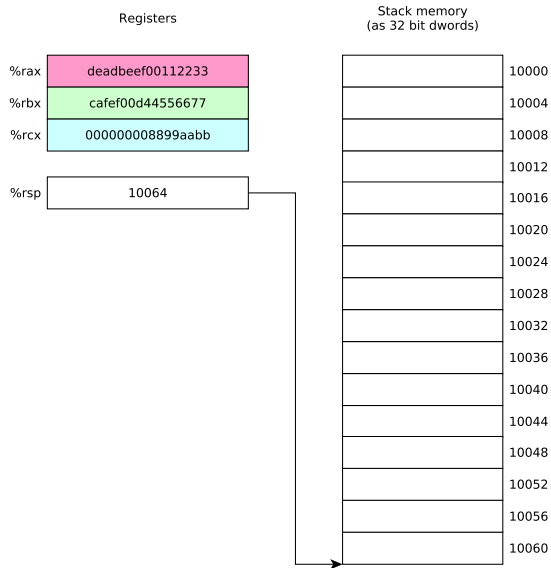
- *Stack pointer* register `%rsp`: contains address of current “top” of stack
  - Important: stack grows towards lower addresses, so top of stack is at lower address than bottom of stack
- *Instruction pointer* register `%rip`: contains code address of next instruction to be updated
  - Control flow changes the value of `%rip`
- Other architectures use the name “program counter” rather than “instruction pointer”, but they’re the same thing

# push and pop

- push: push a data value onto the call stack
  - E.g., `pushq %rax`
    - Decrement `%rsp` by 8
    - Store value in `%rax` at memory location pointed-to by `%rsp`
- pop: pop a data value from the call stack
  - E.g., `popq %rax`
    - Load value at memory location pointed-to by `%rsp` into `%rax`
    - Increment `%rsp` by 8
- `push` and `pop` are amazingly useful for saving and restoring register values
- Various size operands (1, 2, 4, 8 bytes) can be pushed and popped; need to consider alignment

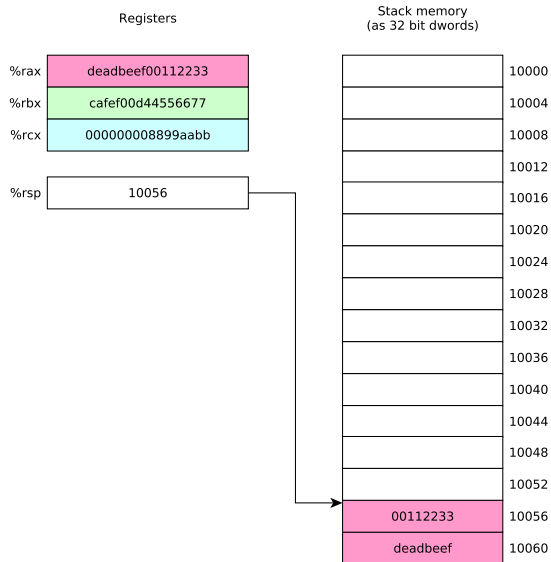


# push and pop



```
pushq %rax
pushq %rbx
pushq %rcx
popq %rbx
popq %rax
popq %rcx
```

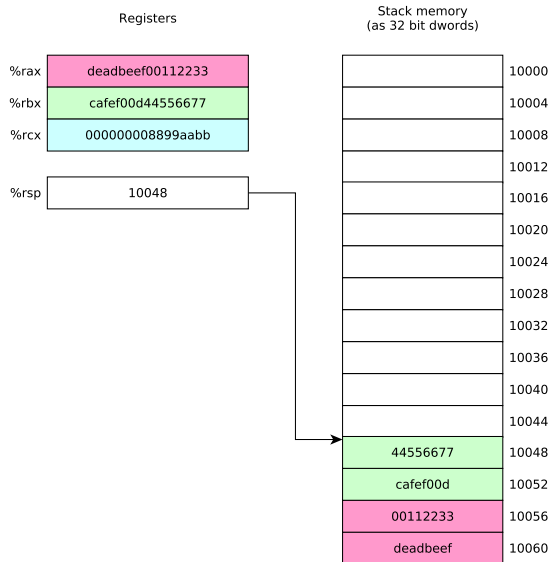
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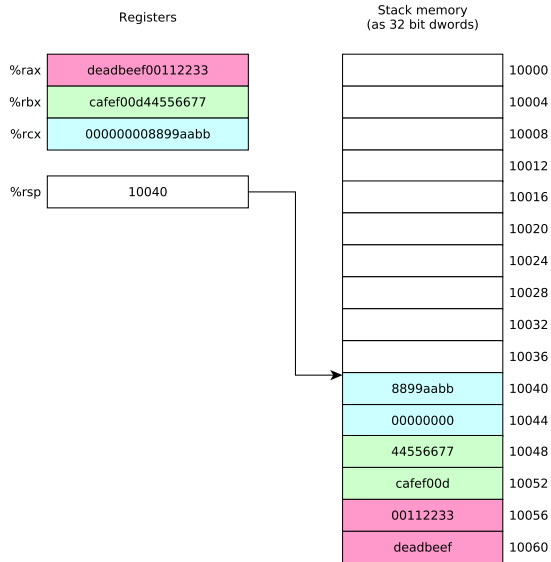


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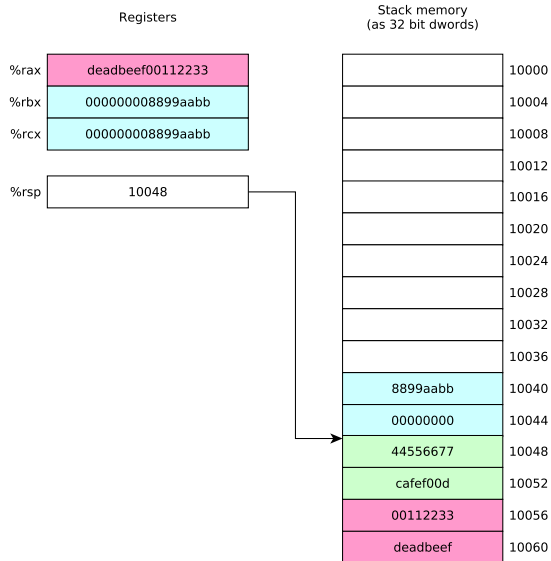
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# push and pop



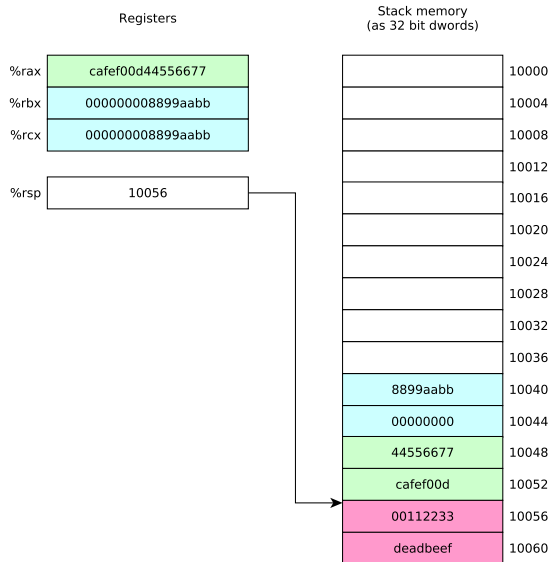
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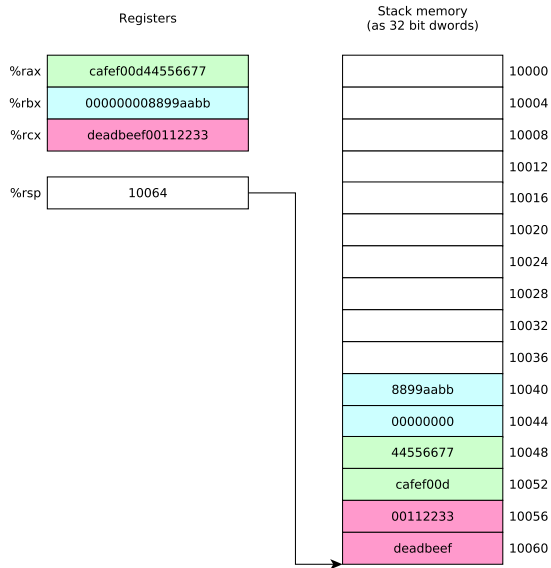
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```

# call and ret

- `call` instruction: calls procedure
  - `%rip` contains address of instruction following `call` instruction
  - Push `%rip` onto stack (as though `pushq %rip` was executed): this is the *return address*
  - Change `%rip` to address of first instruction of called procedure
  - Called procedure starts executing
- `ret` instruction: return from procedure
  - Pop saved return address from stack into `%rip` (as though `popq %rip` was executed)
  - Execution continues at return address



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- This is true of stack-allocated values!
- The Linux x86-64 calling conventions require `%rsp` to be a multiple of 16 at the point of a procedure call (to ensure that 16 byte values can be accessed on the stack if necessary)
- **Issue:** on entry to a procedure,  $\text{\%rsp} \bmod 16 = 8$  because the `call` instruction (which called the procedure) pushed `%rip` (the program counter) onto the stack



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- The Linux `printf` function will segfault if the stack is misaligned





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  - Procedure return value is typically returned in a specific register





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  - They allow your code to interoperate with other code, including library routines and (OS) system calls
- **Always follow the appropriate register use conventions**



# x86-64 Linux register use conventions

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- Callee-saved registers: `%rbx`, `%rbp`, `%r12`, `%r13`, `%r14`, `%r15`



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- *Caller-saved* registers: caller must *not* assume that the procedure call will preserve their value
  - In general any procedure can freely modify them
  - A caller might need to save their contents to memory prior to calling a procedure and restore the value afterwards



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    - Understand that called procedures could modify them!
  - Use callee-saved registers for longer term values that need to persist across procedure calls
    - Use `pushq/popq` to save and restore their values on procedure entry and exit





# Recursive Fibonacci computation

Compute  $n$ th Fibonacci number recursively (warning: exponential-time algorithm!)

The call stack inherently allows recursion: there is nothing special we need to do to make it work

Recall that

$$\text{fib}(0) = 0$$

$$\text{fib}(1) = 1$$

$$\text{For } n > 1, \text{fib}(n) = \text{fib}(n - 2) + \text{fib}(n - 1)$$



# Recursive Fibonacci function (see fibRec.S for full program)

```
fib:
    cmpl $2, %edi          /* check base case */
    jae .LrecursiveCase   /* if n>=2, do recursive case */
    movl %edi, %eax        /* base case, just return n */
    ret

.LrecursiveCase:
    /* recursive case */
    pushq %r12             /* preserve value of %r12 */
    movl %edi, %r12d        /* save n in %r12 */
    subl $2, %edi           /* compute n-2 */
    call fib               /* compute fib(n-2) */
    movl %r12d, %edi        /* put saved n in %edi */
    subl $1, %edi           /* compute n-1 */
    movl %eax, %r12d        /* save fib(n-2) in %r12 */
    call fib               /* compute fib(n-1) */
    addl %r12d, %eax        /* return fib(n-2)+fib(n-1) */
    popq %r12              /* restore value of %r12 */
    ret                    /* done */
```



## Running the program (with $N=9$ )

```
$ gcc -c -g -no-pie -o fibRec.o fibRec.S
$ gcc -no-pie -o fibRec fibRec.o
$ ./fibRec
fib(9) = 34
```



# Clicker quiz!

Clicker quiz omitted from public slides



# Stack memory allocation

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- Could use heap allocation (i.e., `malloc`, `free`)
  - Has overhead due to bookkeeping, locking
- The call stack is an ideal place to allocate storage for local variables



# Stack allocation

- Stack allocation of storage is simple:
  - To allocate  $n$  bytes, subtract  $n$  from `%rsp`
    - Updated `%rsp` is a pointer to the beginning of the allocated memory
  - To deallocate  $n$  bytes, add  $n$  to `%rsp`
- Complication: instructions such as `push` and `pop` change `%rsp`
- Solution: use the *frame pointer* register `%rbp` to keep track of allocated memory area



# Using the frame pointer

On entry to procedure:

```
pushq %rbp
movq %rsp, %rbp
subq $N, %rsp
```

Before returning from procedure:

```
addq $N, %rsp
popq %rbp
```

%rbp points to a memory location *just above* a block of  $N$  bytes allocated in the current stack frame. Note that

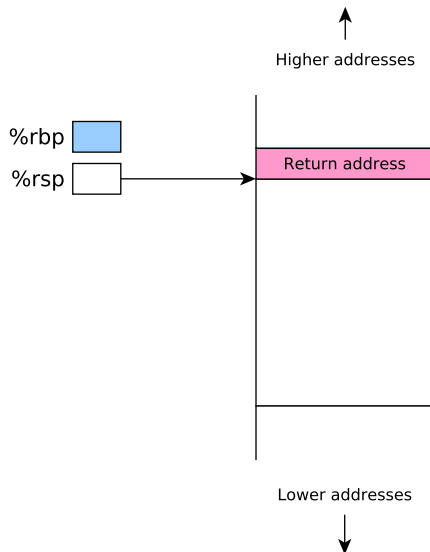
- $N$  should be a multiple of 16 to ensure correct stack alignment
- The function will access memory locations in the allocated block using *negative* offsets from %rbp





# Before allocating space in stack frame

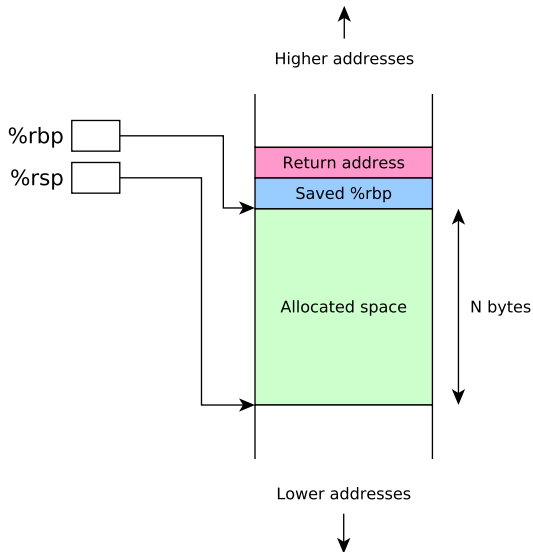
```
--> pushq %rbp  
    movq %rsp, %rbp  
    subq $N, %rsp
```



# After allocating space in stack frame

```
pushq %rbp  
movq %rsp, %rbp  
subq $N, %rsp
```

-->



# Putting it all together

- Let's examine a simple program which
  - Reads two 64 bit integer values from user
  - Computes their sum using a function
  - Prints out the sum
- Calling `scanf` to read input requires variables in which to store input values: we'll allocate them on the stack



# addLongs, C version

```
#include <stdio.h>

long addLongs(long a, long b);

int main(void) {
    long x, y, sum;
    printf("Enter two integers: ");
    scanf("%ld %ld", &x, &y);
    sum = addLongs(x, y);
    printf("Sum is %ld\n", sum);
}

long addLongs(long a, long b) {
    return a + b;
}
```



# addLongs, assembly version

```
.section .rodata
sPromptMsg: .string "Enter two integers: "
sInputFmt:  .string "%ld %ld"
sResultMsg: .string "Sum is %ld\n"
```

```
.section .text
.globl main
.align 16

main:
    pushq %rbp
    movq %rsp, %rbp
    subq $16, %rsp

    movl $0, %eax
    movq $sPromptMsg, %rdi
    call printf

    movl $0, %eax
    movq $sInputFmt, %rdi
    leaq -16(%rbp), %rsi
    leaq -8(%rbp), %rdx
    call scanf
```

```
    movq -16(%rbp), %rdi
    movq -8(%rbp), %rsi
    call addLongs

    movq $sResultMsg, %rdi
    movq %rax, %rsi
    call printf

    addq $16, %rsp
    popq %rbp
    ret

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addLongs:
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    pushq %rbp          <-- save orig value of %rbp
    movq %rsp, %rbp
    subq $16, %rsp

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    movq $sPromptMsg, %rdi
    call printf

    movl $0, %eax
    movq $sInputFmt, %rdi
    leaq -16(%rbp), %rsi
    leaq -8(%rbp), %rdx
    call scanf

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    movq -8(%rbp), %rsi
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    movq $sResultMsg, %rdi
    movq %rax, %rsi
    call printf

    addq $16, %rsp
    popq %rbp
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    call printf

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    subq $16, %rsp

    movl $0, %eax
    movq $sPromptMsg, %rdi
    call printf

    movl $0, %eax
    movq $sInputFmt, %rdi
    leaq -16(%rbp), %rsi <-- pass address of 1st var
    leaq -8(%rbp), %rdx
    call scanf

    movq -16(%rbp), %rdi
    movq -8(%rbp), %rsi
    call addLongs

    movq $sResultMsg, %rdi
    movq %rax, %rsi
    call printf

    addq $16, %rsp
    popq %rbp
    ret

    .align 16
addLongs:
    movq %rdi, %rax
    addq %rsi, %rax
    ret
```



# addLongs, assembly version

```
.section .rodata
sPromptMsg: .string "Enter two integers: "
sInputFmt:  .string "%ld %ld"
sResultMsg: .string "Sum is %ld\n"

.section .text
.globl main
.align 16
main:
    pushq %rbp
    movq %rsp, %rbp
    subq $16, %rsp

    movl $0, %eax
    movq $sPromptMsg, %rdi
    call printf

    movl $0, %eax
    movq $sInputFmt, %rdi
    leaq -16(%rbp), %rsi
    leaq -8(%rbp), %rdx    <-- pass address of 2nd var
    call scanf

    movq -16(%rbp), %rdi
    movq -8(%rbp), %rsi
    call addLongs

    movq $sResultMsg, %rdi
    movq %rax, %rsi
    call printf

    addq $16, %rsp
    popq %rbp
    ret

    .align 16
addLongs:
    movq %rdi, %rax
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main:
    pushq %rbp
    movq %rsp, %rbp
    subq $16, %rsp

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    movq $sPromptMsg, %rdi
    call printf

    movl $0, %eax
    movq $sInputFmt, %rdi
    leaq -16(%rbp), %rsi
    leaq -8(%rbp), %rdx
    call scanf
```

```
movq -16(%rbp), %rdi  <-- pass value of 1st var
movq -8(%rbp), %rsi
call addLongs
```

```
movq $sResultMsg, %rdi
movq %rax, %rsi
call printf
```

```
addq $16, %rsp
popq %rbp
ret
```

```
.align 16
addLongs:
    movq %rdi, %rax
    addq %rsi, %rax
    ret
```



# addLongs, assembly version

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    pushq %rbp
    movq %rsp, %rbp
    subq $16, %rsp

    movl $0, %eax
    movq $sPromptMsg, %rdi
    call printf

    movl $0, %eax
    movq $sInputFmt, %rdi
    leaq -16(%rbp), %rsi
    leaq -8(%rbp), %rdx
    call scanf
```

```
    movq -16(%rbp), %rdi
    movq -8(%rbp), %rsi    <-- pass value of 2nd var
    call addLongs
```

```
    movq $sResultMsg, %rdi
    movq %rax, %rsi
    call printf
```

```
    addq $16, %rsp
    popq %rbp
    ret
```

```
.align 16
addLongs:
    movq %rdi, %rax
    addq %rsi, %rax
    ret
```



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```

```
.section .text
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    pushq %rbp
    movq %rsp, %rbp
    subq $16, %rsp

    movl $0, %eax
    movq $sPromptMsg, %rdi
    call printf

    movl $0, %eax
    movq $sInputFmt, %rdi
    leaq -16(%rbp), %rsi
    leaq -8(%rbp), %rdx
    call scanf
```

```
    movq -16(%rbp), %rdi
    movq -8(%rbp), %rsi
    call addLongs
```

```
    movq $sResultMsg, %rdi
    movq %rax, %rsi
    call printf
```

```
    addq $16, %rsp  <-- deallocate alloc'ed area
    popq %rbp
    ret
```

```
.align 16
addLongs:
    movq %rdi, %rax
    addq %rsi, %rax
    ret
```



# addLongs, assembly version

```
.section .rodata
sPromptMsg: .string "Enter two integers: "
sInputFmt:  .string "%ld %ld"
sResultMsg: .string "Sum is %ld\n"
```

```
.section .text
.globl main
.align 16

main:
    pushq %rbp
    movq %rsp, %rbp
    subq $16, %rsp

    movl $0, %eax
    movq $sPromptMsg, %rdi
    call printf

    movl $0, %eax
    movq $sInputFmt, %rdi
    leaq -16(%rbp), %rsi
    leaq -8(%rbp), %rdx
    call scanf
```

```
    movq -16(%rbp), %rdi
    movq -8(%rbp), %rsi
    call addLongs
```

```
    movq $sResultMsg, %rdi
    movq %rax, %rsi
    call printf
```

```
    addq $16, %rsp
    popq %rbp      <-- restore orig value of %rbp
    ret
```

```
.align 16
addLongs:
    movq %rdi, %rax
    addq %rsi, %rax
    ret
```



# addLongs, assembly version

```
.section .rodata
sPromptMsg: .string "Enter two integers: "
sInputFmt: .string "%ld %ld"
sResultMsg: .string "Sum is %ld\n"
```

```
.section .text
.globl main
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main:
    pushq %rbp
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    movl $0, %eax
    movq $sPromptMsg, %rdi
    call printf

    movl $0, %eax
    movq $sInputFmt, %rdi
    leaq -16(%rbp), %rsi
    leaq -8(%rbp), %rdx
    call scanf
```

```
    movq -16(%rbp), %rdi
    movq -8(%rbp), %rsi
    call addLongs
```

```
    movq $sResultMsg, %rdi
    movq %rax, %rsi
    call printf
```

```
    addq $16, %rsp
    popq %rbp
    ret
```

```
.align 16
```

```
addLongs: <-- does not use stack, ignore alignment :-P
    movq %rdi, %rax
    addq %rsi, %rax
    ret
```



# Running the program

```
$ gcc -c -no-pie -o addLongs.o addLongs.S
$ gcc -no-pie -o addLongs addLongs.o
$ ./addLongs
Enter two integers: 2 3
Sum is 5
```





# Running the program in gdb

```
$ gdb addLongs
...output omitted...
(gdb) break addLongs.S:28
Breakpoint 1 at 0x401172: file addLongs.S, line 28.
(gdb) run
Starting program: /home/daveho/.../src/control2/addLongs
Enter two integers: 3 4

Breakpoint 1, main () at addLongs.S:28
28                movq -16(%rbp), %rdi        /* pass first value */
(gdb) print *(long *)($rbp-16)
$1 = 3
(gdb) print *(long *)($rbp-8)
$2 = 4
```



# Running the program in gdb

```
$ gdb addLongs
...output omitted...
(gdb) break addLongs.S:28  <-- set breakpoint just after scanf returns
Breakpoint 1 at 0x401172: file addLongs.S, line 28.
(gdb) run
Starting program: /home/daveho/.../src/control2/addLongs
Enter two integers: 3 4

Breakpoint 1, main () at addLongs.S:28
28          movq -16(%rbp), %rdi      /* pass first value */
(gdb) print *(long *)($rbp-16)
$1 = 3
(gdb) print *(long *)($rbp-8)
$2 = 4
```



# Running the program in gdb

```
$ gdb addLongs
...output omitted...
(gdb) break addLongs.S:28
Breakpoint 1 at 0x401172: file addLongs.S, line 28.
(gdb) run                                <-- start the program
Starting program: /home/daveho/.../src/control2/addLongs
Enter two integers: 3 4

Breakpoint 1, main () at addLongs.S:28
28          movq -16(%rbp), %rdi          /* pass first value */
(gdb) print *(long *)($rbp-16)
$1 = 3
(gdb) print *(long *)($rbp-8)
$2 = 4
```



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...output omitted...
(gdb) break addLongs.S:28
Breakpoint 1 at 0x401172: file addLongs.S, line 28.
(gdb) run
Starting program: /home/daveho/.../src/control2/addLongs
Enter two integers: 3 4    <-- enter input values

Breakpoint 1, main () at addLongs.S:28
28          movq -16(%rbp), %rdi    /* pass first value */
(gdb) print *(long *)($rbp-16)
$1 = 3
(gdb) print *(long *)($rbp-8)
$2 = 4
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Starting program: /home/daveho/.../src/control2/addLongs
Enter two integers: 3 4

Breakpoint 1, main () at addLongs.S:28
28          movq -16(%rbp), %rdi          /* pass first value */
(gdb) print *(long *)($rbp-16)  <-- print first input value at -16(%rbp)
$1 = 3
(gdb) print *(long *)($rbp-8)
$2 = 4
```



# Running the program in gdb

```
$ gdb addLongs
...output omitted...
(gdb) break addLongs.S:28
Breakpoint 1 at 0x401172: file addLongs.S, line 28.
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Starting program: /home/daveho/.../src/control2/addLongs
Enter two integers: 3 4

Breakpoint 1, main () at addLongs.S:28
28                movq -16(%rbp), %rdi        /* pass first value */
(gdb) print *(long *)($rbp-16)
$1 = 3
(gdb) print *(long *)($rbp-8)    <-- print second input value at -8(%rbp)
$2 = 4
```



# Running the program in gdb

```
$ gdb addLongs
...output omitted...
(gdb) break addLongs.S:28
Breakpoint 1 at 0x401172: file addLongs.S, line 28.
(gdb) run
Starting program: /home/daveho/.../src/control2/addLongs
Enter two integers: 3 4

Breakpoint 1, main () at addLongs.S:28
28          movq -16(%rbp), %rdi      /* pass first value */
(gdb) print *(long *)($rbp-16)
$1 = 3
(gdb) print *(long *)($rbp-8)
$2 = 4
```



# Acknowledgements

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